

51.15 Notions of distance for graphs

Recalled the length $l(\Gamma)$ of a path graph Γ is # of constituent edges.

Definition 1 The distance $d(x, y)$ between vertices x and y in a graph G is given by

$$d(x, y) = \min_{\substack{\text{paths } \Gamma(x \rightarrow y) \\ x \text{ to } y}} (l(\Gamma)) \quad (1);$$

in the event of there being no such path, $d(x, y) := \infty$.

Remark Paths of maximal length are sometimes useful too, c.f. Veblen's Thm in §1.16

Definition 2 The eccentricity $e(v)$ of a vertex v is

$$e(v) = \max_{w \in V(G)} (d(v, w)) = \max_{w \in V(G)} \left(\min_{\text{paths } \Gamma(x \rightarrow y)} (l(\Gamma)) \right)$$

Definition 3 The radius $r(G)$ of a graph G is

$$r(G) = \min_{v \in V(G)} (e(v)),$$

$$\text{i.e. } r(G) = \min_{v \in V(G)} \left(\max_{w \in V(G)} (d(v, w)) \right)$$

$$\text{i.e. } r(G) = \min_{v \in V(G)} \left(\max_{w \in V(G)} \left(\min_{\text{paths } \Gamma(x \rightarrow y)} (l(\Gamma)) \right) \right)$$

cf. Hausdorff dist
& Gromov-Wasserstein dist
in metric space
thing for comparability
long sum
meetings
of max/mins or
inf/sups...

Definition 4 A vertex v is central if $e(v) = r(G)$.

The centre $C(G)$ of a graph $= \{v \mid e(v) = r(G)\}$;
the set of central vertices.

Definition 5 The diameter $d(G)$ of a graph is the opposite extreme to the radius:

$$\begin{aligned} d(G) &= \max_{v \in V(G)} (e(v)) \\ &= \max_{v \in V(G)} \left(\max_{w \in V(G)} (d(v, w)) \right) \\ &= \max_{v \in V(G)} \left(\max_{w \in V(G)} \left(\min_{\text{paths } \Gamma(x \rightarrow y)} (l(\Gamma)) \right) \right) \end{aligned}$$

Definition 6

Lemma $r(G) \stackrel{\textcircled{1}}{\leq} d(G) \stackrel{\textcircled{2}}{\leq} 2r(G).$

Proof

(1) follows from the basic inequality $\min \leq \max$.

For (2) : $\exists x \in G$ st $d(x, v) \leq r(G)$ $\textcircled{1}$ by the minimality of radius.
This holds for any v .

In particular then it holds for a y, z pair which realize the diameter
 $d(y, z) = d(G)$. $\textcircled{*}$

$$\text{so } (1) \Rightarrow \begin{cases} d(y, v) \leq r(G) & \textcircled{2} \\ d(x, v) \leq r(G) & \textcircled{3} \end{cases}$$

But $d(\cdot, \cdot)$ is a distance, so the Δ inequality holds for it :

$$d(y, z) \leq d(x, v) + d(v, x) \leq r(G) + r(G) = 2r(G)$$

$$\textcircled{*} \parallel \text{distance is symmetric} \quad \overbrace{d(x, v)}^{d(G)}$$

□

Definition 6 The girth $g(G)$ of a graph is

$$g(G) = \min_{\text{cycles } C \subseteq G} (l(C))$$

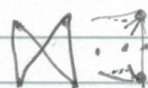
For technical convenience, if G contains no cycles, $g(G) := \infty$
i.e. G (acyclic graph) $:= \infty$
 G (forest)

Examples $G(C_n) = n$

$$G(K_n) = 3$$



$$G(K_{n,n}) = 4$$



$$G(n\text{-pan}) = G(C_n \# P_1) = l(C_n) = n$$

$$G(n\text{-fan}) = 3$$



$$G(n\text{-wheel}) = 3$$



Remark Girth feeds into the concept of cage, in which there is an Ex in §1.16.